

1 **Affine automata verifiers**

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7 We initiate the study of the verification power of Affine finite automata (AfA) as a part of Arthur-Merlin (AM) proof systems. We  
8 show that every unary language is verified by a real-valued AfA verifier. Then, we focus on the verifiers restricted to have only  
9 integer-valued or rational-valued transitions. We observe that rational-valued verifiers can be simulated by integer-valued verifiers,  
10 and their protocols can be simulated in nondeterministic polynomial time. We show that this bound is tight by presenting an AfA  
11 verifier for NP-complete problem SUBSETSUM. We also show that AfAs can verify certain non-affine and non-stochastic unary  
12 languages.  
13

14 CCS Concepts: • **Theory of computation** → **Formal languages and automata theory**.

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16 Additional Key Words and Phrases: affine automata, interactive proof systems , Arthur-Merlin games, unary languages, subset-sum  
17 problem, NP.  
18

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23  
24 Affine finite automata (AfAs) are quantum-like generalization of probabilistic finite automata (PFAs) mimicking quantum  
25 interference and having the capability of “making measurement” based on  $\ell_1$ -norm (called weighting). The computation  
26 of an AfA is linear, but the weighting operators may be non-linear.  
27

28 AfAs was formally defined in [3], and it was shown that they are more powerful than PFAs and quantum finite  
29 automata (QFAs) in bounded error and unbounded error settings, but their nondeterministic version is equivalent to  
30 nondeterministic QFAs. Since then, AfAs and their different generalizations (e.g., OBDDs and using counters) have  
31 been investigated in a series of work [4–6].  
32

33 We consider AfAs as part of Arthur-Merlin (AM) proof systems and investigate their verification power. We show that  
34 every unary language can be verified by a real-valued AfA verifier. Then, we focus on the verifiers with integer-valued  
35 or rational-valued transitions. We show how to simulate rational-valued verifiers by integer-valued ones and how to  
36 simulate their protocols in nondeterministic polynomial time. We present an AfA verifier for NP-complete problem  
37 SUBSETSUM. We also show that AfAs can verify certain non-affine and non-stochastic unary languages. In our protocols,  
38 we use similar verification strategies and encoding techniques previously used for two-way QFAs in [7, 8].  
39

40 We define AM(AfA) as the class of languages verifiable by bounded-error Arthur-Merlin system having realtime  
41 affine finite verifiers. In [2], Arthur-Merlin systems with probabilistic finite automata verifier is defined as an automata  
42 having both nondeterministic and probabilistic states. We follow the same framework here. We indeed give the ability  
43 of making nondeterministic transitions to the model of a finite automaton with classical and affine states.  
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53 The list of standard complexity classes mentioned in the paper:

54	REG	:	regular languages
55	L	:	logarithmic space
56	P	:	polynomial time
57	NP	:	nondeterministic polynomial time
58	SPACE(n)	:	linear space
59	PSPACE	:	polynomial space
60	NEXP	:	nondeterministic exponential space
61	ALL	:	all languages

62 On unary languages, for the real-valued verifiers, we show that AfAs and 2QCFA have the same verification power,  
63 where 2QCFA is the two-way QFA model defined in [1]:

64  $UALL = UAM(2QCFA) = UAM(AfA)$ , where AfAs are realtime machines but 2QCFA run in exponential expected  
65 time.

66 On unary languages, for the rational-valued verifiers, we know that

$$67 \text{UREG} = UAM_{\mathbb{Q}}(\text{PFA}) \subseteq \begin{matrix} UAM_{\mathbb{Q}}(\text{QFA}) \subseteq UAM(\text{QFA}) \\ UAM_{\mathbb{Q}}(2\text{PFA}) \subseteq UAM(2\text{PFA}) \end{matrix},$$

68 where it is open whether the inclusions are strict, and we show that  $UPOLY(P) \in UAM_{\mathbb{Q}}(AfA)$  and so we have  
69  $UREG \subseteq UAM_{\mathbb{Q}}(AfA)$ .

70 On non-unary languages, for the rational-valued verifiers, we give an upper bound for  $AM_{\mathbb{Q}}(AfA)$ , and so we have

$$71 AM_{\mathbb{Q}}(AfA) = AM_{\mathbb{Z}}(AfA) \subseteq NP \cap SPACE(n) \subseteq AM_{\mathbb{Q}}(2QCFA),$$

72 where 2QCFA run in double-exponential expected time. Our bound is tight since we show that  $SUBSETSUM \in AM_{\mathbb{Z}}(AfA)$ .

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